

Computer Networks

Basics

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Agenda

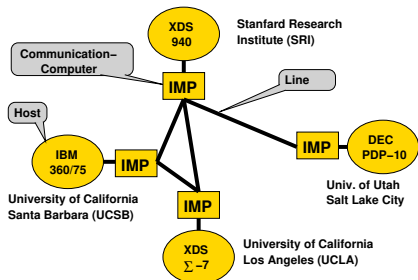
- Historical background
- Components and Terms
- Reference Models
- Topologies

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- Topologies

The ARPANET

- 1957 Foundation of the **Advanced Research Projects Agency (ARPA)** by the US Dept of Defense (DoD) in response to *Sputnik*
- 1962 The idea of the '**Internet**' as 'tool to create critical mass of intellectual resources' (Licklider, Taylor)
- 1967 Plan for the **ARPANET** was published
Main architects: *Vinton Cerf, Bob Kahn*
- 1969 First **Request for Comments (RFC)** and first **functioning network**, rented 50 kBit/sec lines, Interface Message Processors by BBN



Graphic by courtesy of Prof. Dr. Roland Kaiser, Hochschule RheinMain

First Internet Protocols

- 1972 First public demo (remote login) using the Network Control Protocol (NCP)
main use: terminal sessions, file transfer, Electronic Mail
- 1974 Basics of TCP/IP written on paper by Cerf/Kahn (IP=Internet Protocol, TCP=Transmission Control Protocol), standardization in the following years
- 1982 Transition towards IP version 4 (IPv4)¹
- from 1983 Dissemination of TCP/IP due to Berkeley UNIX 4.2 BSD, source code publicly available



Author: Gorthmog

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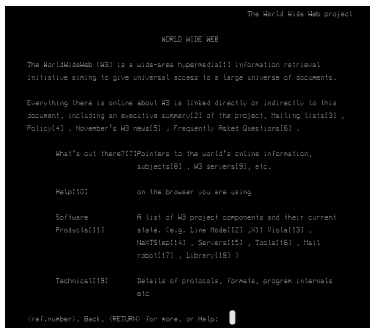


²deprecated, but still widely used

²<https://creativecommons.org/licenses/by-sa/4.0/deed.en>

Standardization

- 1986 The **Internet Engineering Task Force (IETF)** is founded as an open standardization organization
- 1989 Foundation of **RIPE (Réseaux IP Européens)** as a forum for administrative and technical coordination of Internet development
- 1990 Proposal of a hypertext project at CERN in Geneva by Tim Berners-Lee and Robert Cailliau: cradle of the **world wide web**
- 1995 The specification of **IPv6** (as a successor of IPv4) is published by the IETF



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³<http://line-mode.cern.ch/www/hypertext/WWW/TheProject.html>

Global Success

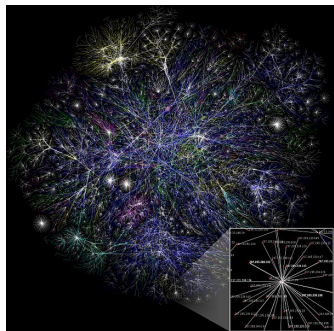
- 1996 First **search engines** with a site-scoring algorithm, e.g., Google search
- 1998 Start of the **dot-com** boom
- 2004 Start of **Web 2.0** brought up blogs and RSS as well as services like Facebook or Twitter
- 2007 Apple's iPhone and Android started the "**Mobile Revolution**"
- 2008 Rise of the **Internet of Things (IoT)**



<https://bit.ly/3JDGA3w>

Internet growth

- Amount of **AS** (Autonomous Systems, admin. routing domain)
 - **Doubling** every five years (currently, more than 100,000)
 - Stable core
 - Major growth at the fringe
- Traffic rate
 - Growth rate of about **26% per year** estimated
- Users
 - 2021: **two third** of the world population is "online"⁴
 - More than **doubled** during the last **ten years**⁴
 - Strongest growth outside the EU, Japan, and USA⁴



https://en.wikipedia.org/wiki/File:Internet_map_1024.jpg

⁴Source: <https://www.internetworldstats.com/stats.htm>

Agenda

■ Historical background

■ Components and Terms

■ Reference Models

■ Topologies

Purpose of Computer Networks

The general task of a computer network is to enable communication among the participants.

- Resource sharing
 - ⇒ assign different tasks to different computers
 - ⇒ avoid bottlenecks
- Resource pooling
 - ⇒ combine the resources and functionalities of multiple machines
- Resource balancing
 - ⇒ increase the availability of the services by redundancy

Required Components to set up a Computer Network

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 - Rules that specify, how computers can communicate

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Some of the technologies, concepts, and terms are used in a different contexts. For example, network services communicating on one host or connected peripheral devices within one computer device.

Network Services

- A **network service** provides resources to other devices in the network
- Distinguished by their role:
 - Server** Provides a network service
 - Client** Uses (consumes) a network service
- If each communication partner is server and client both, the participants are called **peers** (\implies Peer-to-Peer networks)
- The terms server, client and peer typically refer only to network services and not to hardware
 - Reason: It is common that client applications also run at *servers*

Transmission Media

Different transmission media exists to setup a computer network.

1 Guided transmission media

- **Copper cable:** Data is transferred as electrical impulses
- **Fiber-optic cable:** Data is transferred as light impulses

2 Wireless transmission

- Wireless transmission can be realized **directed** and **undirected**
- Directed transmission can base on the following technologies:
 - **Radio technology:** Data is transferred as electromagnetic waves (radio waves) in the radio frequency spectrum (e.g., directed WLAN and satellite internet access)
 - **Infrared:** Data is transferred as electromagnetic waves in the spectral range (e.g., IrDA)
 - **Laser:** Data is transferred as light impulses via Laser Bridge
- Undirected wireless transmission is always based on radio technology (e.g., WLAN, cellular networks, terrestrial broadcasting and satellite broadcasting)

Protocols

- A **protocol** is the set of all previously made **agreements** between communication partners
 - These agreements include:
 - Rules for **connection establishment** and **termination**
 - Method of **synchronization** between sender and receiver (if any)
 - Measures for the **detection and treatment of transmission errors**
 - Definition of **valid messages** (vocabulary)
 - **Format and encoding** of messages
- Protocols specify. . .
 - the **syntax** (= format of valid messages)
 - the **semantics** (= vocabulary and meaning of valid messages)



Computer Networks distinguished by their Dimension (1/3)

- Depending on the dimension, different groups of computer networks are distinguished
- **Personal Area Network (PAN) or Body Area Network (BAN)**
 - Network of small mobile devices, such as smart phones
 - Technologies: USB, FireWire, WLAN, Bluetooth, IrDA
 - Major dimension: Few meters
- **Local Area Network (LAN)**
 - Local network
 - Range covers an apartment, building, company site or university campus
 - Major dimension: 500-1000 m
 - Concrete values depend on the transmission medium used and when using wireless networks, also the environment and the transmission power
 - Technologies: Ethernet, Wireless LAN (WLAN), Token Ring (outdated)

Computer Networks distinguished by their Dimension (2/3)

■ Metropolitan Area Network (MAN)

- Connects LANs
- Range covers a city or agglomeration area
- Major dimension: 100 km
- Technologies: Fiber-optic cables, WiMAX (IEEE 802.16)
 - Fiber-optic cables are used because of lesser attenuation (signal weakening) and higher data transmission rates

■ Wide Area Network (WAN)

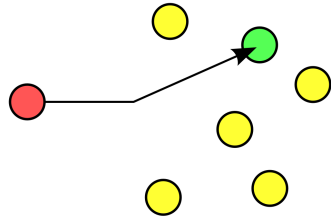
- Connects several networks
- Range covers a large geographic area inside a country or continent
- Major dimension: 1000 km
- Technologies: Ethernet (10 Gbit/s), Asynchronous Transfer Mode (ATM)

Communication Modes

- **Synchronous** ("Rendez-Vous")
 - Sender and receiver needs to be present at the same time
 - May require to **wait** for the other side to become ready
 - For example, phone calls or video conference
- **Asynchronous**
 - Sender and receiver may act independently from each other
 - Requires **buffering**
 - For example, instant messaging or E-Mail

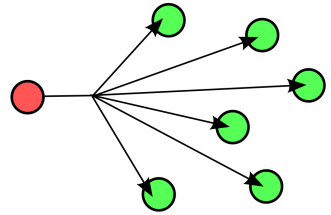
Unicast and Broadcast

Unicast One-to-one communication, i.e., one host sends information to *exactly one* other host



Source: public domain

Broadcast One-to-all communication, i.e., one host sends information to *all* other hosts in the network

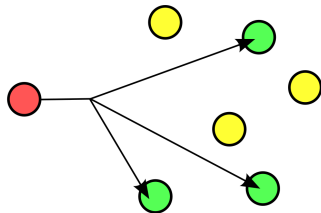


Source: public domain

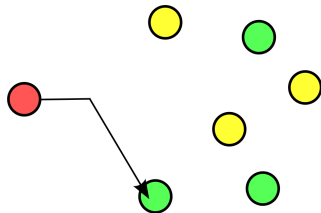
Group Communication: Multicast and Anycast

Multicast Group communication, i.e., one host sends information to all hosts in a given *group*

Anycast One-to-any communication, i.e., one hosts sends information to *one host* in a given *group*



Source: public domain



Source: public domain

Connection-Orientation

Network services may operate *connection-oriented* or *connectionless*.

connection-oriented the service operates **stateful**

- comprises three phases: connection establishment, data transfer, and connection termination
- a virtual path between the involved hosts is established
- sequent data is exchanged between the same hosts
- typically used for reliable services

connectionless the service operates **stateless**

- no path between the involved hosts is established
- typically used for low latency services

Directional Dependence (Anisotropy) of Data Transmission

Given a **communication channel** with two (or more) **endpoints**:

■ Simplex

- Only one side of the channel can send data → the channel can be used in only one direction
- Examples: Radio, TV, Pager

■ Duplex (Full-duplex)

- Both sides of the channel are allowed to send → the channel can be used in both directions simultaneously
- Examples: Phone, Networks with twisted pair cables because they provide separate wires for send and receive

■ Half-duplex

- Both sides of the channel can send, but not simultaneously → the channel can only be used in one direction at a time
- Examples:
 - Networks with fiber-optic cables or coaxial cables, because there exists just a single line to sending and receiving
 - Wireless networks with just a single channel

Bandwidth, Throughput and Goodput

- Main factors, influencing the performance of a computer network:
 - **Bandwidth** (→ throughput)
 - **Latency (delay)**
- The **bandwidth** specifies how many bits can be transmitted within a period via the network
 - If a network has a bandwidth of 1 Mbit/s, one million bits can be transmitted per second **in the ideal case**
 - Thus, a bit has a *width* of $1 \mu\text{s}$
 - If the bandwidth is doubled, the number of bits that can be transmitted per second double, too
 - **Throughput** is the actual achieved data rate (\Rightarrow the bandwidth defines its upper bound)
 - **Goodput** is the actual rate of data the user benefits from

Latency

- The **latency** of a network is the time, a message needs to travel from one end of the network to the most distant end

Latency = Propagation delay + Transmission delay + Waiting time

$$\text{Propagation delay} = \frac{\text{Distance}}{\text{Speed of light} * \text{Velocity factor}}$$

- Distance: Length of the network connection
- Speed of light: 299, 792, 458 m/s
- Velocity factor: Vacuum = 1, twisted pair cables = 0.6, optical fiber = 0.67, coaxial cables = 0.77

$$\text{Transmission delay} = \frac{\text{Message size}}{\text{Bandwidth}}$$

Transmission delay = 0, if the message consists only of a single bit

- Waiting times are caused by network devices (e.g., Switches)
 - They need to cache received data first before forwarding it
 - ⇒ Waiting time = 0, if the network connection between sender and destination is just a single line or a single channel

Source: Larry L. Peterson, Bruce S. Davie. Computernetzwerke. dpunkt (2008)

Bandwidth-Delay Product

- Calculates the **volume of a network connection**
 - Signals cannot be transmitted with infinite speed via the transmission media
 - The propagation speed is in any event limited by the speed of light and it depends on the velocity factor of the transmission medium
 - The product of bandwidth and delay (latency) corresponds to the maximum number of bits that can reside inside the line between sender and receiver
- Example: A network with 100 Mbit/s bandwidth, and 10 ms latency

$$100,000,000 \text{ Bits/s} \times 0.01 \text{ s} = 1,000,000 \text{ Bits}$$

- There are a maximum number of 1,000,000 Bits inside the network line
 - This is equivalent to 125,000 Bytes (approx. 123 kB)

How does a Computer Network work?

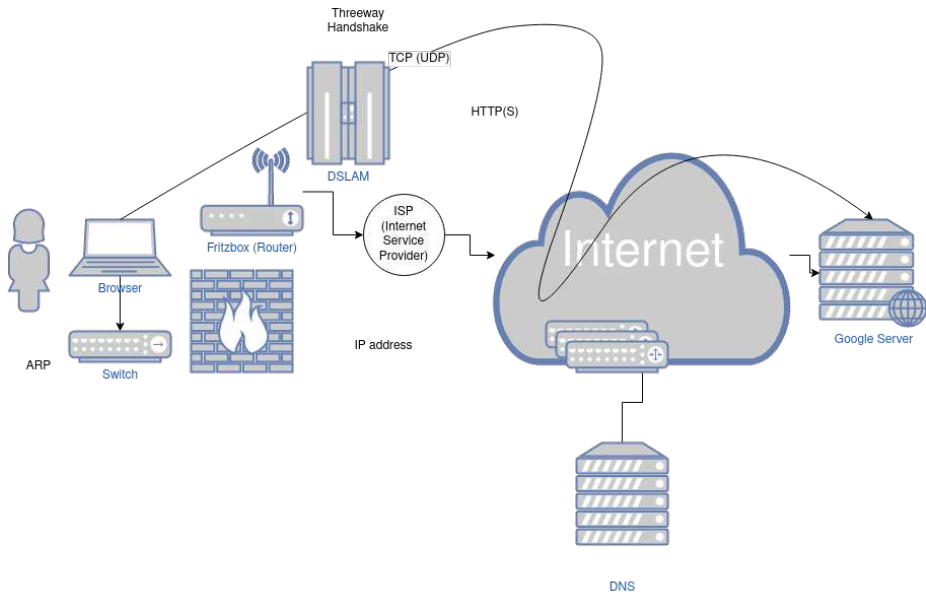
You need information about someone/something:

- What do you do?

How does a Computer Network work?

- You need information about someone/something:
- What do you do?
 - Which problems are to solve?

The Big Picture



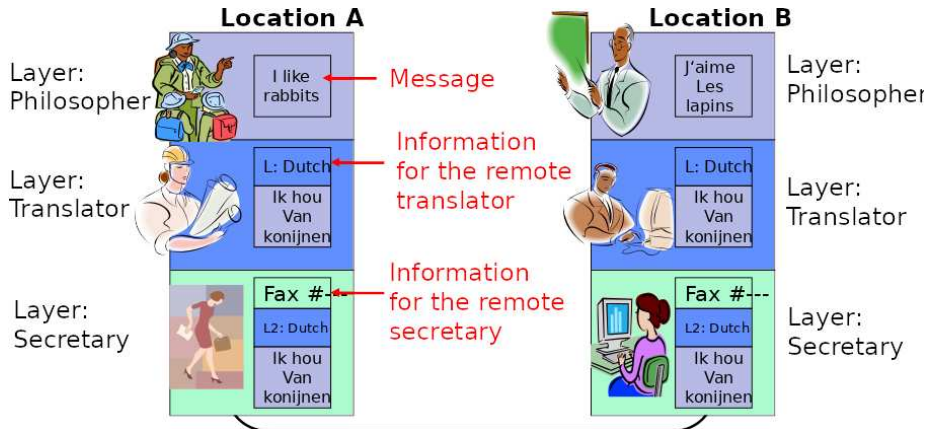
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Reference Models

- **Reference models** are used to describe computer networks independently of concrete technologies
- Such a reference model consists of several layers
- Each **layer** addresses a particular aspect of communication and offers interfaces to the neighboring layer
- Each layer defines their own protocols that define syntax and semantics of parts of a transmitted message (e.g., header and trailer)
- These message parts are **encapsulated**
- Because each layer is complete in itself, single protocols can be modified or replaced without affecting all aspects of communication
- The most popular reference models are...
 - the **TCP/IP reference model**,
 - the **ISO/OSI reference model**, and
 - the **hybrid reference model**

"Philosopher-Translator-Secretary"-Architecture



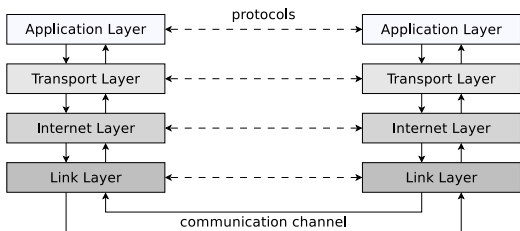
Graphic by courtesy of Prof. Dr. Thomas C. Schmidt, HAW Hamburg

TCP/IP Reference Model or DoD Model

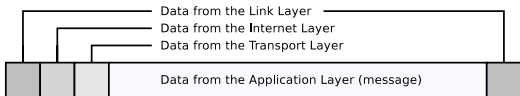
- Developed from 1970 onwards by the Department of Defense (DoD) in the Arpanet project
- Divides the required functionality to realize communication into 4 layers
- For each layer, it is specified, what functionality it provides
 - These requirements are implemented by communication protocols
 - Concrete implementation is not specified and can be implemented in different ways
 - Therefore, for each of the 4 layers, multiple protocols exist

Number	Layer TCP/IP (RFC 1122)	Layer DoD (RFC 871)	Protocols (Examples)
4	Application Layer	Process Layer	HTTP, FTP, SMTP, POP3, DNS, SSH, Telnet
3	Transport Layer	Host-to-Host Layer	TCP, UDP
2	Internet Layer		IPv4, IPv6, IPX
1	Link Layer	Network Interface Layer	Ethernet, WLAN, ATM, FDDI, PPP, Token Ring

TCP/IP Reference Model – Message Structure

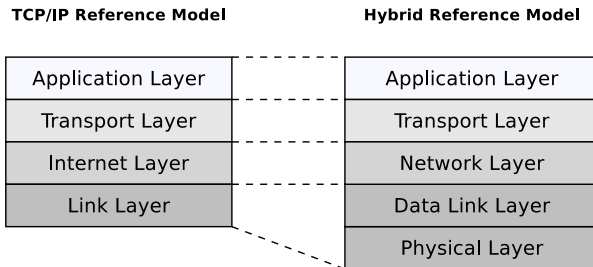


- Each layer adds additional information as **header** to the message
 - Some protocols (e.g., Ethernet) add in the link layer not only a header but also a **trailer** at the end of the message
 - The receiver analyzes the header (and trailer) on the same layer



Hybrid Reference Model

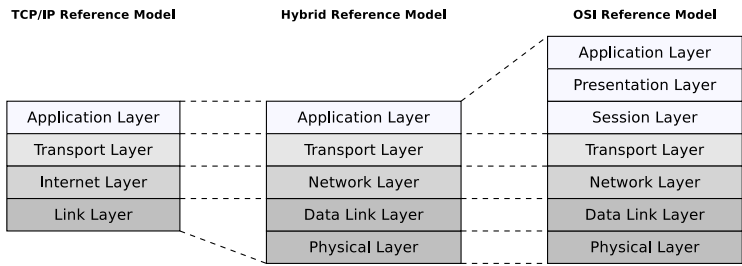
- The TCP/IP reference model is often presented in the literature (e.g., by Andrew S. Tanenbaum) as a 5-layer model
 - Reason: It makes sense to split the **Link Layer** into 2 layers, because they have different tasks
- This model is an extension of the TCP/IP model and is called **hybrid reference model**



We will mostly follow the hybrid reference model

OSI Reference Model

- Some years after the TCP/IP reference model (1970s), the OSI (Open Systems Interconnection) reference model was developed from 1979 onwards
- 1983: Standardized by the Intern. Organization for Standardization (ISO)
- In contrast to the hybrid reference model, two additional layers are placed below the Application and above the Transport Layer



OSI Model Concepts

Central concepts of the OSI model are:

Services Define what the layer does, i.e., its semantics

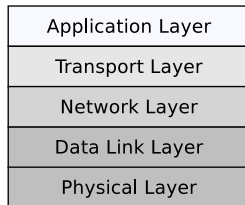
Interfaces Define how to access it

Protocols Describe how the layer is implemented

Physical Layer I

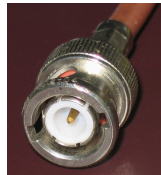
- **Transmits the ones and zeros**
 - **Physical connection** to the network
 - Conversion of data into **signals**
- Protocol and transmission medium specify among others:
 - How is the information encoded on the transmission medium?
 - Can transmission take place simultaneously in both directions?

Hybrid Reference Model



Physical Layer II

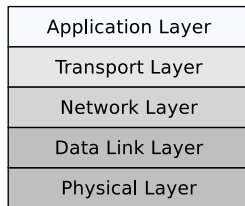
- **At sender site:** Signals are modulated onto the medium
- **At receiver site:** Signals are demodulated from the medium
- **Devices: Repeater, Hub (Multiport Repeater)**



Data Link Layer I

- Ensures **error-free** data exchange of **frames** between devices in physical networks
 - Handles transmission errors with **checksums**
 - Controls the access to the transmission medium (e.g., via CSMA/CD or CSMA/CA)
- Specifies physical network addresses (**MAC addresses**)

Hybrid Reference Model



Data Link Layer II

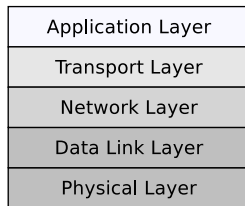
- **At sender site:** Packs the Network Layer packets into frames and transmits them (in a reliable way) via a physical network from one device to another
- **At receiver site:** Identifies frames in the bit stream from the Physical Layer
- Devices: **Bridges, Layer-2-Switches** (Multiport Bridges), **WIFI APs**, and **Modems** connect physical networks



Network Layer I

- Forwards **packets** between logical networks (over physical networks)
 - For this *internetworking*, the network layer defines **logical addresses** (most commonly **IP addresses**)
 - Each IP packet is **routed** independently to its destination (→ connectionless)

Hybrid Reference Model



Network Layer II

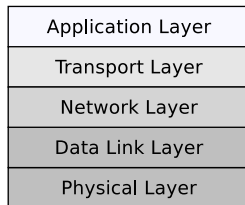
- **At sender site:** Packs the segments of the Transport Layer in packets
- **At receiver site:** Unpacks the packets in the frames from the Data Link Layer
- **Routers and Layer-3-Switches** connect logical networks
- Usually the connectionless Internet Protocol (IP) is used
 - Other protocols (e.g., IPX) have been replaced by IP



Transport Layer I

- Transports **segments** between processes on different devices via so-called end-to-end protocols
- Transport protocols implement different forms of communication
 - **Connectionless** communication, typically UDP (User Datagram Protocol) in TCP/IP networks
 - **Connection-oriented** communication, typically TCP (Transport Control Protocol) in TCP/IP networks

Hybrid Reference Model



Transport Layer II

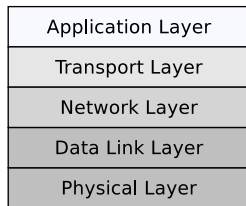
- **At sender site:** Packs the data of the Application Layer into segments
- **At receiver site:** Unpacks the segments inside the packets from the network layer
- Addresses processes with **port numbers**

Combination of TCP/IP = de facto standard for computer networks

Application Layer

- Contains all protocols, that interact with the **application programs** (e.g., browser or email program)
- Here is the actual **payload** (e.g., HTML pages or emails), formatted according to the used application protocol
- Some Application Layer protocols: HTTP, FTP, SMTP, POP3, DNS, SSH, Telnet

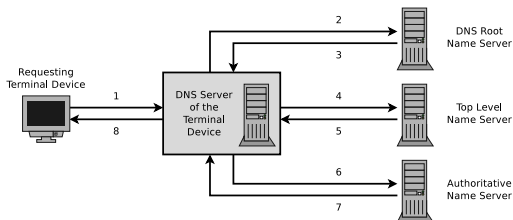
Hybrid Reference Model



wikipedia.org (CC0)



pixabay.com (CC0)



OSI only: Session Layer

- **Controls the dialogues** (connections) between processes
- Provides the following services
 - **checkpointing** (and recovery)
 - **authentication**
 - **authorization**
- Relevant protocols of the Session Layer are H.245, L2TP, PAP, and SOCKS
- Session Layer services are commonly used for RPCs (cf. lecture *Distributed Systems*)

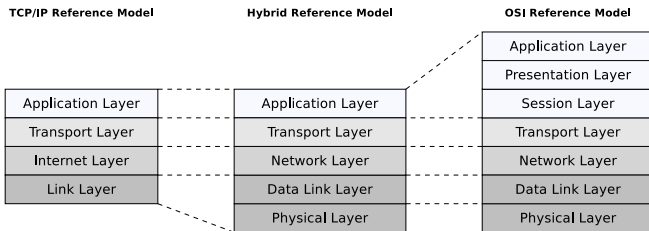
OSI only: Presentation Layer

- Contains rules for setting the **format (presentation)** of messages
 - The sender can notify the receiver that a message has a specific **format** (e.g., ASCII) to make conversion happen, which is perhaps necessary
 - Data records can be specified here with fields (e.g., name, student ID number...)
 - **Data types and their length** can be defined here
 - **Compression and encryption** could be implemented by this layer

The Presentation Layer is seldom used in practice, because all tasks intended to this layer are fulfilled by Application Layer protocols today

Reference Models – Summary

- Conclusion: The hybrid reference model illustrates the functioning of computer networks in a realistic way
 - It distinguishes between the Physical Layer and Data Link Layer
 - This is useful, because the objectives differ a lot
 - It does not subdivide the Application Layer
 - This is less helpful and often not realized in practice
 - Functionalities, which are intended for Session Layer and Presentation Layer, are provided by Transport or Application Layer protocols and services

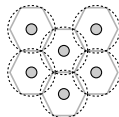
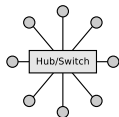
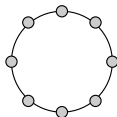
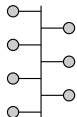


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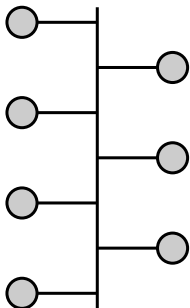
Topologies of Computer Networks

- The topology of a computer network. . .
 - determines how the communication partners are connected with each other
 - affects its reliability a lot
- The structure of large-scale networks is often a combination of different topologies
- Physical and logical topology may differ
 - **Physical topology**: Describes the wiring
 - **Logical topology**: Describes the flow of data between the terminal devices
- Topologies are graphically represented with nodes and edges



Bus Network

- All terminal devices are connected via a **shared communication medium** – the **bus**

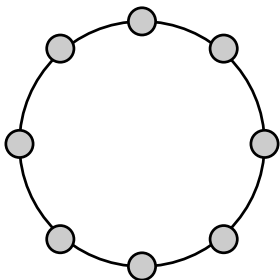


- No active components between the terminal devices and the shared communication cable
 - If a node fails, it does not affect the network itself
- **Advantage:** Cheap to implement
 - In the past, Hubs and Switches have been expensive
- **Drawback:** Shared communication cable fails
 - ⇒ Complete network fails
- Only a single node can send data at each point in time
 - ⇒ otherwise, collisions will occur
 - A media access control method like CSMA/CD is required

- Examples:

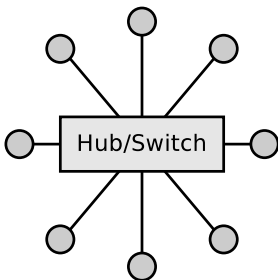
- (original) Ethernet, CAN, I2C

Ring Network



- Connects node to node
- All data is transferred from nodes to nodes until the destination is reached
- Disruption of a single link \implies network failure
- Each node is also a repeater, which amplifies the signal
 - For that reason, large-sized rings (transmission medium dependent) are possible
 - Maximum ring length for Token Ring: 800 m
- Examples:
 - Token Ring (**logical**): 4-16 Mbps
 - Fiber Distributed Data Interface (FDDI): 100-1000 Mbps
 - FDDI implements 2 rings
 - One is a secondary backup, in case the primary ring fails

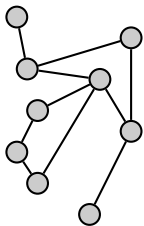
Star Network



- All nodes are connected directly with a central component (Hub or Switch)
 - Failure of the central component leads to a failure of the network itself
 - The central component can be implemented in a redundant way
 - Failure of a node do not cause a failure of the network itself
 - **Advantages:** Expandability and stability
- Examples:
- (modern) Ethernet
 - Token Ring (**physical**): 4-16 Mbps
 - Fibre Channel (storage networks): 2-16 Gbps
 - InfiniBand (cluster): 10-40 Gbps

Mesh Network

- Each node is connected with one or more other nodes
 - In a **fully connected mesh network**, the nodes are all connected to each other
- If nodes or connections fail, communication inside the network is typically still possible because the frames are redirected



- **Advantages:** Failure safe (depends on the degree)
- **Drawbacks:** Cabling effort and energy consumption
- Additional challenge: complexity to find the best way from sender to receiver (cf. *Travelling salesman problem*)
- Examples:
 - Logical topology between Routers
 - Ad-hoc (wireless) networks

Tree Network

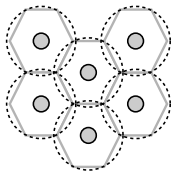
- A dedicated **root** node exist with one or more edges
 - Every edge leads to a **leaf** node or to the root of another tree
- Several star topology networks are hierarchically connected
- **Advantages:**
 - Failure of a terminal device (leaf node) has no consequences
 - Good expandability and long distances are possible
 - Well suited for searching and sorting algorithms
- **Drawbacks:**
 - When a node fails, the complete (sub-)tree behind is no longer accessible
 - In a large tree, the root may become a bottleneck because the communication from one half of the tree to the other half always needs to pass the root



- **Example:**
 - Connecting Hubs or Switches via an uplink port

Cellular Network

- Implemented by wireless networks
- **Cell**: Area where the nodes can communicate with the base station
- **Advantage**: Failure of nodes do not affect the network itself
- **Drawback**: Maximum dimension is limited by the number of base stations and their positions



- Only one nodes can send data at each point in time
⇒ otherwise, collisions will occur
 - A media access control method like CSMA/CA is required
- Examples:
 - Wireless LAN = WiFi (*IEEE 802.11*)
 - Global System for Mobile Communications (*GSM*)

Current Situation

- Today, Ethernet (1-10 Gbit/s) with Switches (\implies **star topology**) is the standard for wired LAN
- Connecting Hubs and Switches implements a **tree topology**, if there are no loops in the cabling
- **Cell topology** is the standard for wireless networks
- **Mesh topology** is one possible use case of wireless networks and it is the logical topology between routers
- **Bus and ring topologies** are no longer used for new computer network infrastructures
 - 10BASE2 (Thin Ethernet) and 10BASE5 (Thick Ethernet) are outdated since the mid/end-1990s
 - May 2004: IBM sells his complete Token Ring product lineup

You should now be able to answer the following questions:

- What is a Computer Network and what are its objectives?
- What is the difference between bandwidth, throughput, and latency?
- What is a reference model and what do their difference layers represent?

