

# Distributed Systems

## Network Services

Prof. Dr. Oliver Hahm

Frankfurt University of Applied Sciences  
Faculty 2: Computer Science and Engineering  
[oliver.hahm@fb2.fra-uas.de](mailto:oliver.hahm@fb2.fra-uas.de)  
<https://teaching.dahahm.de>

29.04.2024

# Agenda

- Internet Protocols

- Network Performance

# Agenda

■ Internet Protocols

■ Network Performance

- Which functions needs to Be implemented By a computer network?

- Which functions needs to Be implemented By a computer network?
- How can we handle all these functionalities?

- Which functions needs to Be implemented By a computer network?
- How can we handle all these functionalities?

■ → Layers

- Which functions needs to Be implemented By a computer network?
- How can we handle all these functionalities?

■ → Layers

- What is a layer?

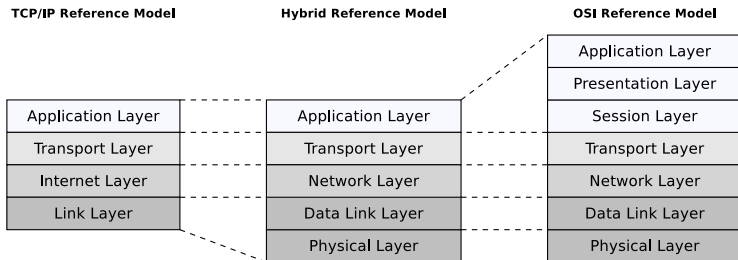
# OSI Reference Model

Central concepts of the **OSI** model are:

**Services** Define what the layer does, i.e., its semantics

**Interfaces** Define how to access it

**Protocols** Describe how the layer is implemented





What are the tasks, devices, and protocols for each layer?

# Tasks of Layer 1

- Layer 1: Physical Layer (PHY)
  - Definition of mechanical/electrical/optical specifications and procedures required for the Bit transmission on the carrier.
  - The Physical Layer can be subdivided into
    - 1 *PMD Sublayer*: Physical Media Dependent (mechanical/electrical/optical) specifications
    - 2 *PHY Sublayer*: Physical (Signals ↔ bits: en/de-coding)
  - Transmits the ones and zeros

# Tasks of Layer 2

- **Layer 2: Data Link Layer (DL):**
  - Provides mechanisms how two nodes on the same physical network guarantee a safe communication.
  - For this purpose, the information (= bits) are grouped together in *frames* and are complemented by checksum as final part of the frame. The checksum allows an error detection (and potentially correction).
  - The Data Link Layer can be subdivided into two sublayers as well:
    - 1 **MAC Sublayer: Media Access Control** defines the access to the physical carrier and how nodes are addressed
    - 2 **LLC Sublayer Logical Link Control** provides an additional safety and control functionality

# Tasks of Layers 3 and 4

## ■ Layer 3: Network Layer

- The task of this layer is to transmit the data in *packets* between the (network) nodes.
- Forwards *packets* between logical networks (over physical networks)
- Thus, layer 3 can be understood as packet exchange layer.

## ■ Layer 4: Transport Layer

- The transport layer provides a (potentially *reliable*) virtual *end-to-end connection* for the data transport between the end nodes
- Acts as a *multiplexer* between the various processes on the hosts via **ports**
- Transport protocols can implement **connection-oriented** or **connectionless** communication

# Tasks of Layers 5 and 6

## ■ Layer 5: *Session Layer*

- This is the lowest application-specific layer and is responsible to raise, maintain, and gracefully terminate communication relationships between the end nodes: a *Session*.
- It's particular scope is to provide dialog-functionalities among the communication partners, in particular to allow a synchronization between the involved communication processes.

## ■ Layer 6: *Presentation Layer*

- The conversion and translation of different data representations (e.g. character set families like [ASCII](#) and [EBCDIC](#)) to a common format prior of sending, is main task of the Presentation Layer.
- This layer may include functionalities which allow compression of data, conversion, and encryption.
- Known presentation schemes here were [ASN.1](#) (*Abstract Syntax Notation No. 1*) and [XDR](#) (*eXternal Data Representation*); however [XML](#) (*eXtensible Markup Language*) is now mostly used instead.

# Tasks of layers 7

## ■ Layer 7: Application Layer

- Contains all protocols, that interact with the **application programs** (e.g., browser or email program)
- Here is the actual **payload** (e.g., HTML pages or emails), formatted according to the used application protocol
- Comprises the biggest variety of protocols

# Main Characteristics of the TCP/IP Architecture

- **Connectionless best-effort** protocol (**IP**) on the network layer
- **Packet switching** via network nodes
- Static and dynamic **routing**
- Transport protocol with reliable **end-to-end transmissions** (**TCP**)

What does packet switching mean?



# Circuit versus Packet switching

## ■ Circuit switching

- A sustained (virtual) **connection** is present between both communication partners (end nodes) as long as the data transmission lasts
- End nodes must be attached to a specific exchange node
- **Drawback:** The communication breaks down, in case the connection fails

# Circuit versus Packet switching

## ■ Circuit switching

- A sustained (virtual) **connection** is present between both communication partners (end nodes) as long as the data transmission lasts
- End nodes must be attached to a specific exchange node
- **Drawback:** The communication breaks down, in case the connection fails

## ■ Packet switching

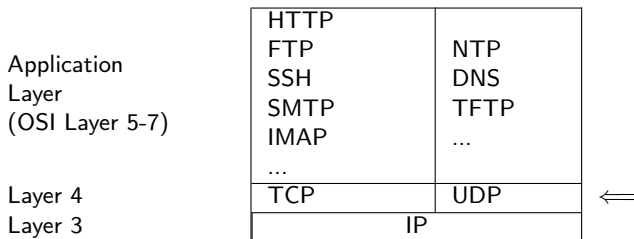
- The data to be exchanged will be encapsulated in packets as complete **information units**
- These packets are dropped on the network and exchanged between the communication partners
  - Packets do include an address-information about the sender and the recipient node
  - Packets can be buffered on the transmission path
- **Drawback:** Some per-packet overhead for long-lasting connections

# Essential Properties of IP

- **Connectionless** protocol
- **best-effort** transport of individual messages  
(Datagram, =Packet)
- Addressing of hosts via 32 bit (IPv4) or 128 bit (IPv6) **IP addresses**
- *Fragmentation* if necessary (optional in IPv6)
- IPv4 contains a checksum for the header, but not the payload; IPv6 does not contain any checksum
- IPv4 packets may contain optional header fields, IPv6 uses extension headers
- The **lifetime** (*TTL*) of a packet in the network is limited

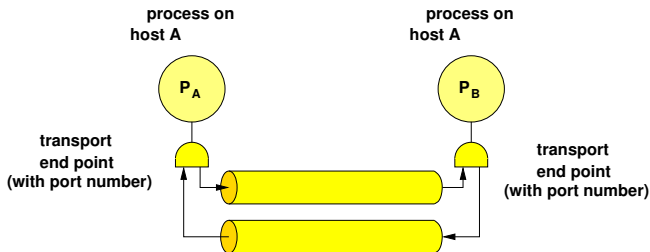
# Transport Layer Protocols

- Tasks of the **Transmission Control Protocols (TCP)**:
  - Reliable bidirectional point-to-point transport of a bytestream between two hosts on the endpoints
- Tasks of the **User Datagram Protocols (UDP)**:
  - Best-effort datagram service of IP layer is accessible for the processes on the endpoints
- **Classification:**



# TCP Communication Model

- connection oriented
- virtual, **bidirectional**, full-duplex capable connection between endpoints (used by their processes)
- Addressing of transport layer endpoints via 16 bit **port numbers** (in addition to the node's IP addresses)
- **Bytestream** oriented, not blockwise
- A single packet containing a chunk of the bytestream is called **segment**



# Main Characteristics of TCP

- **Reliable** transmission due to ...
  - Sequence numbers
  - Checksum calculation (same algorithm as IPv4)
  - Reception receipts (acknowledgements) and timeouts
  - Retry after timeout
- Preserved **order**
- *Sliding window* principle for **flow control**
- *AIMD* principle for **congestion control**

# Main Characteristics of TCP

- **Reliable** transmission due to ...
  - Sequence numbers
  - Checksum calculation (same algorithm as IPv4)
  - Reception receipts (acknowledgements) and timeouts
  - Retry after timeout
- Preserved **order**
- *Sliding window* principle for **flow control**
- *AIMD* principle for **congestion control**
  - **AIMD** := Additive Increase/Multiplicative Decrease

# Essential Properties of UDP

- **Connectionless** protocol
- Addressing of the user via 16 bit **port numbers**
- *best-effort* transmission of **datagrams** (individual messages), IP services from the network layer are made accessible for application processes → 1 : 1 mapping
- Multicast/Broadcast capable (1:n communication), direct application of multicast capable networks like, for instance, Ethernet
- Integrity check via an optional checksum
- No further guarantees:
  - No receipts or other guarantees, i.e., datagrams can get lost, arrive in a different order, or getting duplicated
  - No flow control, i.e., on the receiver site datagrams may get discarded in case of full or missing buffers
- Well suited for the implementation of simple request/response protocols



For which type of application  
would you use TCP and for  
which UDP?

- Which properties of a distributed system are affected by the Link Layer?
- What is the impact of packet switching (e.g., compared to circuit switching) on the Network Layer on a distributed system?
- What are the criteria to select the Transport Layer protocol when designing a distributed application?

# Agenda

■ Internet Protocols

■ Network Performance

- How can we measure the network performance?

- How can we measure the network performance?
- What does affect the network performance?

# Effecting the Network Performance

- Data and messages, which are transmitted over networks may be **lost** and/or **corrupted**:
  - Insufficient quality of the transmission carrier
  - External distortion impacts (e.g., because of electromagnetic fields)
- The communication protocol has to deal with these cases and has to provide:
  - **Error/Fault Control**:  
Identification and compensation for transmission errors/failures
  - **Flow Control**:  
Adaptive means to adjust the amount of data to be send w.r.t. the recipient's (announced) capacity
  - **Congestion Control**:  
Additional means, to reduce the potential lost of data (packets) on the network

# Failure Causes

## Signal Transmission Errors

During the transmission of bit sequences on the physical layer errors may occur

They are typically caused by...

- **Signal deformation**
  - Attenuation of the transmission medium
- **Noise**
  - Thermal or electronic noise
- **Crosstalk**
  - Interference by neighboring channels
  - Capacitive coupling increases with increasing frequency
- **Short-time disturbances**
  - Cosmic radiation
  - Defective or insufficient insulation

## Error Types

Burst errors are more common than single bit errors

## Typical BER values

POTS	$2 * 10^{-4}$
Radio link	$10^{-3} - 10^{-4}$
Ethernet	$10^{-9} - 10^{-10}$
Fiber	$10^{-10} - 10^{-12}$

The LLC sublayer tries to **detect** and **handle** bit errors that occur during signal

# Checksum

## Checksum

The checksum is calculated by a pre-defined algorithm for a block of data. They are typically used for the verification of the data integrity.

- For error detection, the sender attaches a **checksum** at each frame
- The receiver can now detect erroneous frames and **discard** them
- Possible checksums:
  - **Parity-check codes**
  - The polynomial code – **Cyclic Redundancy Checks (CRCs)**



# Error Control

- In order to detect transmission errors on the upper layers positive  $A^+$  and negative **acknowledgments**  $A^-$  are feasible.
  - However, acknowledgments can be corrupted or get lost as well
  - The sender has to consider a *deferment period* until the acknowledgment has been finally received
  - In addition, the data blocks (or even the transmitted byte) can be *enumerated* (bookkeeping).
- The sender has to keep the transmitted data in his *sending buffer* until he finally has received the acknowledgment.

# Flow control

- Flow control enables the **adaption of the transmission rate** of the sender with respect to
  - the recipient or
  - any network component which is responsible for the data transfer
- Typical flow control methods:
  - Messages **hold and continue** (XON/XOFF) issued by the recipient also know as Ready-for-Reception/Clear-to-Send (RFR/RTS+CTS),
  - By issuing **credits**
  - **Window mechanism** where the communication partners mutually tell their reception buffer to each other and adjust the data to transmit according to the provided value

# Congestion control

- Any physical network has only a certain **capacity** to transmit only a certain amount of data during a certain time period.
- Congestion occurs when...
  - the recipient buffer in any network component is exaggerated ( $\Rightarrow$  incoming data packets need to be dropped)
  - the sender is required to build up additional send buffers (*queues*) without being able to transmit the data packets on the network
- **congestion avoidance** is task of **congestion control**, since any congestion will impact
  - the data **throughput** and
  - the **transfer latency** (delay)negatively

## Important takeaway messages of this chapter

- The Internet's TCP/IP architecture provides a flexible and generic communication system for many types of higher layer services
- While IP provides a packet switching best-effort service, transport layer protocols can offer additional services
- The network has to manage transmission errors and control the data flow

